

**Exercise set 2****Number Theory****Tuesday SEP 27 2011 at 4 pm. SHARP (!)****in MaD-302**

1. Use Eratosthenes' sieve to find all primes under 200

2. Let  $p \neq 3$  be a prime. Prove that

$$p = 3k + 1 \quad \text{or} \quad p = 3k - 1 \quad \text{for some } k \in \mathbb{N}.$$

3. Prove: if  $p$  is prime and  $a \in \mathbb{Z}$ , then either  $p \mid a$  or  $(a, p) = 1$ .

4. Prove that if  $n$  and  $a$  are natural numbers and  $\sqrt[n]{a} \in \mathbb{Q}$ , then  $\sqrt[n]{a} \in \mathbb{N}$  so for example  $\sqrt[3]{10}$  is irrational.

5. in Euclid's classical proof, a prime outside  $\{p_1, p_2, \dots, p_n\}$  is found by considering prime factors of

$$N_n = p_1 p_2 \cdots p_n + 1$$

. Do this beginning with  $\{2\}$ , next being  $\{2, p_2\}$ , where in fact  $N_2 = 2 + 1 = 3$ , so  $p_2 = 3$  since  $N_2$  happens to be prime. Continue, until

(1) either, you have found 5 odd primes . (or more, if you like)

(2) ir:  $N_p$  is not a prime  $p_n \neq N_n$ .

Idesas? Questions??

6. a) 3, 5 and 7 are a triple of primes:  $p, p + 2, p + 4$  Why are there no others?

b) let  $a, b \in \mathbb{N}$  and  $(a, b) \geq 2$ . prove that the set  $A = \{an + b \mid n = 0, 1, 2, \dots\}$  contains at most one prime.

7. Prove that there is a number  $C > 0$ , such that for all  $k \geq 2$

$$(1) \quad \sum_{p \leq k, p \in \mathbb{P}} \frac{1}{p} \geq \log \log k + C,$$

so the series  $\sum_{p \in \mathbb{P}} \frac{1}{p}$  diverges. You may assume as known (lectures!) that

$$(2) \quad \prod_{p \leq k, p \in \mathbb{P}} \frac{1}{1 - p^{-1}} \geq \sum_n \frac{1}{n} \geq \log k.$$

Take logarithms. Remember how to use them, and notice that

(1)  $-\ln\left(1 - \frac{1}{p}\right) \leq \frac{1}{p} + \frac{1}{p^2}$ , (proof not required today, but easy using series or the fact that  $f(x) = \log(1+x) - x + x^2$  decreases on  $[-\frac{1}{2}, 0]$ )

(2) the series  $\sum_p p^{-2}$  converges.

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8. Calculate (at least some terms of)

- (1)  $E_0 * E_0$
- (2)  $E * E_0$
- (3)  $E_0 * \Omega$
- (4)  $E * N_\alpha$
- (5)  $E * \sigma_{\frac{1}{2}}$
- (6)  $\mu * E * E_0$ .

9. Just read:

Remember : Eukleideen algoritmi luvuille 126 and 35:

$$\begin{aligned} 126 &= 3 \cdot 35 + 21, \\ 35 &= 1 \cdot 21 + 14, \\ 21 &= 1 \cdot 14 + 7, \\ 14 &= 2 \cdot 7. \end{aligned}$$

$s$  and  $t$  are found "backwards":

$$\begin{aligned} (126, 35) &= 7 = 21 - 1 \cdot 14, \\ &= 21 - (35 - 1 \cdot 21), \\ &= (126 - 3 \cdot 35) - (35 - (126 - 3 \cdot 35)), \\ &= 2 \cdot 126 - 7 \cdot 35. \end{aligned}$$

This is clumsy when large numbers on computers. Better:

Let  $\ell, q_i, r_i$  be like in Eukleideen algoritmi. try to find  $s_i$  and  $t_i$  such that  $s_i r_0 + t_i r_1 = r_i$  for all  $0 \leq i \leq \ell$ .

Assume first, that such numbers exist: Apply tis to indices  $i - 1, i$  and  $i + 1$  and use Eukleideen algoritmin:

$$\begin{aligned} (3) \quad r_{i+1} &= r_{i-1} - q_i r_i = (s_{i-1} r_0 + t_{i-1} r_1) - q_i (s_i r_0 + t_i r_1) \\ &= (s_{i-1} - q_i s_i) r_0 + (t_{i-1} - q_i t_i) r_1. \end{aligned}$$

But  $r_{i+1} = s_{i+1} r_0 + t_{i+1} r_1$ . Choos the coefficients recursively:

$$(4) \quad \begin{aligned} s_{i+1} &= s_{i-1} - q_i s_i, \\ t_{i+1} &= t_{i-1} - q_i t_i. \end{aligned}$$

Then, by (3), if  $s_k r_0 + t_k r_1 = r_k$  for  $k = i - 1$  and  $k = i$  and the coefficients  $s_k$  and  $t_k$  are found by (4) then the equation  $s_k r_0 + t_k r_1 = r_k$  is also satisfied for  $k = i + 1$ . So, it is sufficient to find suitable initiala values. Such are

$$s_0 = 1, \quad t_0 = 0, \quad s_1 = 0, \quad t_1 = 1.$$

In the extended Euclidean algorithm, numbers  $\ell, q_i, r_i \in \mathbb{N}, s_i, t_i \in \mathbb{Z}, 1 \leq i \leq \ell$ , are found such that  $0 \leq r_{i-1} < r_i$ , for  $1 \leq i \leq \ell$ , ja

$$(5) \quad \begin{cases} s_0 = 1, & t_0 = 0 \\ s_1 = 0, & t_1 = 1 \\ r_{i-1} = q_i r_i + r_{i+1} \\ s_{i-1} = q_i s_i + s_{i+1} \\ t_{i-1} = q_i t_i + t_{i+1} \end{cases}$$

Then  $s_i r_0 + t_i r_1 = r_i$  for all  $0 \leq i \leq \ell$  and  $r_\ell = (r_0, r_1)$ .  
 Literature [?, §3.2], [?, §4.5.2].

ESIMERKKI. The previous example in the extended algorithm gives

$i$	$r_i$	$s_i$	$t_i$
0	126	1	0
1	35	0	1
2	21	1	-3
3	14	-1	4
4	7	2	-7
5	0	-5	18

Rivittä  $i = 4$  saadaan  
 $r_\ell = (r_0, r_1) = s_\ell r_0 + t_\ell r_1$ , eli  
 $7 = (126, 35) = 2 \cdot 126 - 7 \cdot 35$ .